

## Performance evaluation of hose bandwidth allocation method using feedback control and class-based queuing for VPNs

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## Outline

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## Background: VPNs

- As a *ubiquitous environment* develops, the importance of data transmission increases.
  - ♦ More secure, easy, and stable
- **Virtual Private Networks (VPNs)** enable us to transmit important data securely.
  - ♦ Many global companies have strong interest in constructing *large-scale VPNs*.
- Subscribers and the amount of traffic will increase continuously.
  - ♦ A *starvation* of network resources will occur.
  - ♦ **Scalable QoS mechanism** will be needed.

## Background: Hose model

- Conventional model versus hose model

Model	Conventional model	Hose model
Contract	<i>Pipes</i> (destination sites)	<i>Hose</i> (aggregation of pipes)
Flexibility	<i>Low</i> (static allocation to pipes)	<i>High</i> (dynamic allocation among pipes)
Utilization	<i>Low</i> (Resources are uselessly needed)	<i>High</i> (Reduce required resources)

- Components of VPN hose model
  - ♦ Provisioning method
    - To allocate long term bandwidth to meet subscriber requirements with the minimum bandwidth consumption
  - ♦ **Bandwidth allocation method**
    - To control bandwidth allocation parameters in response to the traffic changes

## Our objective

- To meet *QoS assurance* and *high utilization*, we produce a **hose bandwidth allocation method**.
  - ♦ Without this method, instantaneous changing traffic cannot be accommodated.
- **Our requirements** in terms of QoS assurance
  - ♦ Proportional fair bandwidth allocation among subscribers
  - ♦ Fair bandwidth allocation among active sites with in the allocated bandwidth for the subscriber
  - ♦ High utilization

*An example of fair bandwidth allocation (%)*

Time		0	1	2	3
Subscriber X	Site X1	0	100	50	25
	Site X2	0	0	0	25
Subscriber Y	Site Y1	0	0	50	50

## Hose bandwidth allocation method: Key idea

- **Key idea**
  - ♦ Integration of feedback-driven traffic control and QoS scheduler installation at ingress router

Ingress routers with QoS scheduler

### Hose bandwidth allocation method: Components

- Two components of our method
  - Feedback-driven traffic control: WPFRA
  - Hierarchical packet scheduler: CBQ
- Weight Proportional Fair Rate Allocation
  - Ingress router: receive control packets
  - Egress router: send control packets
  - Core router: update congestion information
- Class-Based Queueing (CBQ)
  - Set up classes at subscriber and site levels

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### Simulation: Model

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### Simulation: Transitional state

**Our method**  
40 to 70s after run simulation  
Throughput to B2 is not much affected when traffic to A3 is started.  
=>Requirements about fair allocation are satisfied even when another flow arrives.

**Original WPFRA**

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### Conclusion

- We proposed a hose bandwidth allocation method for QoS-enabled VPN services which meets three requirements.
  - Proportional fair bandwidth allocation among subscribers
  - Fair bandwidth allocation among active sites
  - High utilization
- Two components of our method
  - Feedback-based traffic control
  - Modified QoS scheduler at ingress router
- Simulation results showed
  - Our method achieved all three requirements in time progress scenario.
- Future work
  - Stability evaluations in large scale scenario

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